

Lecture 11: Categories, functors & functor morphisms, I

1) Categories.

2) Functors.

Ref: [R], Sections 1.1, 1.3

BONUS: Homotopy category of topological space.

1) Our exposition of Category theory will start w. exploring basic notions: categories, functors & functor morphisms.

Definitions below will have a familiar structure: have data & axioms. E.g. here's a basic algebraic structure.

Definition: a **monoid** is

(Data): a set M equipped w. a multipl'n map $M \times M \rightarrow M$

(axioms): that is associative and has unit, 1 .

For example, a group is exactly a monoid, where all elements are invertible. Every ring is a monoid w.r.t. multiplication.

1.1) Definition of a category.

Definition: A **category**, \mathcal{C} , consists of

(Data):

- a "collection" of **objects**, $Ob(\mathcal{C})$.

- $\forall X, Y \in Ob(\mathcal{C}) \rightsquigarrow$ a set of **morphisms**, $Hom_{\mathcal{C}}(X, Y)$

- $\forall X, Y, Z \in Ob(\mathcal{C})$, a map (of sets) called **composition**

$$Hom_{\mathcal{C}}(X, Y) \times Hom_{\mathcal{C}}(Y, Z) \rightarrow Hom_{\mathcal{C}}(X, Z), (f, g) \mapsto g \circ f$$

(\circ is often omitted)

These satisfy:

(Axioms): i) composition is associative:

$(f \circ g) \circ h = f \circ (g \circ h)$ for $f \in \text{Hom}_e(W, X)$, $g \in \text{Hom}_e(X, Y)$, $h \in \text{Hom}_e(Y, Z)$.

ii) Units: $\forall X \in \text{Ob}(\mathcal{C}) \exists 1_X \in \text{Hom}_e(X, X)$ s.t.

• $f \circ 1_X = f \quad \forall f \in \text{Hom}_e(X, Y)$,

• $1_X \circ g = g \quad \forall g \in \text{Hom}_e(Z, X)$.

1.2) Examples

1) Category of sets, **Sets**: objects = sets, morphisms = maps of sets, composition = composition of maps. Axioms: classical (unit $1_X = \text{id}_X$).

2) Sets w. additional str. structure: objects = sets w. add'l str. structure, morphisms = maps compatible w. this str. structure, composition = comp'n of maps. This includes

a) Category of groups, **Groups**: objects are groups, morphisms = homomorphisms of groups.

b) Category of rings, **Rings**.

c) For a ring A , have categories of A -modules, **A -Mod**, & A -algebras (**A -Alg**), in the latter morphisms = A -linear homomorphisms of rings.

Not all categories have the form in 2:

3a) Let Γ be an oriented graph w. vertices V & edges E .

\leadsto category $\mathcal{C}(\Gamma)$, the path category of Γ .

• Objects = V .

• Morphisms = paths in the graph:



this includes empty paths, one for every vertex.

• Composition: concatenation of paths.

Axioms: associativity is manifest, 1_X = empty path in X .

3b) Note: $\forall X \in \text{Ob}(\mathcal{C}) \Rightarrow \text{Hom}_{\mathcal{C}}(X, X)$ is a monoid w.r.t. \circ

Conversely, every monoid, M , gives a category w. one object, X ,
& $(\text{Hom}_{\mathcal{C}}(X, X), \circ) := M$.

1.3) Remarks:

1) Sometimes, objects in a category form a set (here we say our category is **small**). In general, they form a "class", a notion defined in Set theory. We'll ignore this issue.

2) 1_X is uniquely determined. Moreover, if $f \in \text{Hom}_{\mathcal{C}}(X, Y)$, has a (2-sided) inverse g (i.e. $g \in \text{Hom}_{\mathcal{C}}(Y, X) \mid f \circ g = 1_Y, g \circ f = 1_X$) then g is unique, $f^{-1} := g$. In this case, f is called an **isomorphism**; we say X & Y are **isomorphic** (X & Y behave the same from the point of view of \mathcal{C} , e.g. $Z \in \text{Ob}(\mathcal{C}) \leadsto$

$$\begin{array}{ccc} \text{Hom}_{\mathcal{C}}(Z, X) & \xrightarrow{\sim} & \text{Hom}_{\mathcal{C}}(Z, Y) \\ \psi \downarrow & \longmapsto & f \circ \psi \\ \psi & \longmapsto & f \circ \psi \end{array} \quad (\text{inverse is } \psi' \mapsto f^{-1} \circ \psi').$$

Notation: $X \xrightarrow{f} Y$ means $f \in \text{Hom}_{\mathcal{C}}(X, Y)$.

1.4) Subcategories: \mathcal{C} is a category.

Def'n: (i) By a **subcategory**, \mathcal{C}' , in \mathcal{C} we mean:

(Data) • A subcollection, $\text{Ob}(\mathcal{C}')$, in $\text{Ob}(\mathcal{C})$.

• $\forall X, Y \in \text{Ob}(\mathcal{C}')$, a subset $\text{Hom}_{\mathcal{C}'}(X, Y) \subset \text{Hom}_{\mathcal{C}}(X, Y)$ s.t.

(Axioms) • If $f \in \text{Hom}_{\mathcal{C}'}(X, Y)$, $g \in \text{Hom}_{\mathcal{C}'}(Y, Z) \Rightarrow g \circ f \in \text{Hom}_{\mathcal{C}'}(X, Z)$

• $1_X \in \text{Hom}_{\mathcal{C}'}(X, X) \quad \forall X \in \text{Ob}(\mathcal{C}')$.

(ii) A subcategory \mathcal{C}' in \mathcal{C} is called **full** if $\text{Hom}_{\mathcal{C}'}(X, Y) = \text{Hom}_{\mathcal{C}}(X, Y)$, $\forall X, Y \in \text{Ob}(\mathcal{C}')$.

A subcategory \mathcal{C}' has a natural category str'ure.

Examples:

1) A monoid $M =$ category w. one object

A nonempty subcategory M' in $M =$ a submonoid.

M' is full $\Leftrightarrow M' = M$.

2) $\mathbb{Z}\text{-Mod}$ (a.k.a. category of abelian groups) is a full subcategory in Groups

3) The category of commutative rings, CommRings is a full subcategory in Rings .

1.5) Constructions w. categories.

Definition: For a category, \mathcal{C} , its **opposite category**, \mathcal{C}^{opp} consists of

- the same objects as \mathcal{C} ,
- $\text{Hom}_{\mathcal{C}^{\text{opp}}}(X, Y) := \text{Hom}_{\mathcal{C}}(Y, X)$
- $g \circ^{\text{opp}} f := f \circ g$ ($f \in \text{Hom}_{\mathcal{C}^{\text{opp}}}(X, Y) = \text{Hom}_{\mathcal{C}}(Y, X)$,
 $g \in \text{Hom}_{\mathcal{C}^{\text{opp}}}(Y, Z) = \text{Hom}_{\mathcal{C}}(Z, Y)$).

Definition: For categories $\mathcal{C}_1, \mathcal{C}_2$, their **product** $\mathcal{C}_1 \times \mathcal{C}_2$ is defined by:

- $\text{Ob}(\mathcal{C}_1 \times \mathcal{C}_2) = \text{Ob}(\mathcal{C}_1) \times \text{Ob}(\mathcal{C}_2)$
- $\text{Hom}_{\mathcal{C}_1 \times \mathcal{C}_2}((X_1, X_2), (Y_1, Y_2)) = \text{Hom}_{\mathcal{C}_1}(X_1, Y_1) \times \text{Hom}_{\mathcal{C}_2}(X_2, Y_2)$
- composition is componentwise.

Rem: for usual categories we care about (Groups, Rings, A -Mod), the opposite cat'y essentially has no independent meaning, except: $\mathcal{C} = \text{CommRings}$, where \mathcal{C}^{opp} is the category of affine schemes, which is of crucial importance for Algebraic geometry.

2) **Functors:** Motto: a relation between a category and a functor is analogous to a relation between a group and a

homomorphism.

Let \mathcal{C}, \mathcal{D} be categories.

Definition: A **functor** $F: \mathcal{C} \rightarrow \mathcal{D}$ is

(Data)

- an assignment $X \mapsto F(X): \text{Ob}(\mathcal{C}) \rightarrow \text{Ob}(\mathcal{D})$.
- $\forall X, Y \in \text{Ob}(\mathcal{C})$, a map $\text{Hom}_{\mathcal{C}}(X, Y) \rightarrow \text{Hom}_{\mathcal{D}}(F(X), F(Y))$
 $f \mapsto F(f)$

(Axioms) - compatibility between compositions & units

- $\forall f \in \text{Hom}_{\mathcal{C}}(X, Y), g \in \text{Hom}_{\mathcal{C}}(Y, Z) \Rightarrow F(g \circ f) = F(g) \circ F(f)$
equality in $\text{Hom}_{\mathcal{D}}(F(X), F(Z))$.
- $F(1_X) = 1_{F(X)} \forall X \in \text{Ob}(\mathcal{C})$

Example: Let \mathcal{C}, \mathcal{D} be categories w. single object corresponding to monoids M, N . Then a functor $\mathcal{C} \rightarrow \mathcal{D}$ is the same thing as a monoid homomorphism.

Remarks:

- Have the identity functor $\text{Id}_{\mathcal{C}}: \mathcal{C} \rightarrow \mathcal{C}$

- For functors $F: \mathcal{C} \rightarrow \mathcal{D}, G: \mathcal{D} \rightarrow \mathcal{E}$ can take the composition $G \circ F: \mathcal{C} \rightarrow \mathcal{E}$ ($(G \circ F)(X) = G(F(X))$), it's a functor.

- A functor $F: \mathcal{C} \rightarrow \mathcal{D}$ is the same thing as a functor $\mathcal{C}^{\text{opp}} \rightarrow \mathcal{D}^{\text{opp}}$

More examples:

1) Let \mathcal{C}' be a subcategory in \mathcal{C} . Then have inclusion functor $\mathcal{C}' \hookrightarrow \mathcal{C}$ sending objects/morphisms in \mathcal{C}' to the same objects/morphisms now in \mathcal{C} ; axioms are clear.

2) Forgetful functors: forget part of a structure

2a) For: Groups \rightarrow Sets;

On objects: For(G) = G viewed as a set.

On morphisms: For(f) = f , viewed as a map of sets.

Axioms: clear.

BONUS: homotopy category of topological spaces.

B1) Equivalence on morphisms.

Let \mathcal{C} be a category. Suppose that $\forall X, Y \in \text{Ob}(\mathcal{C})$, the set $\text{Hom}_{\mathcal{C}}(X, Y)$ is endowed with an equivalence relation \sim s.t.

(1) If $g, g' \in \text{Hom}_{\mathcal{C}}(Y, Z)$ are equivalent & $f \in \text{Hom}_{\mathcal{C}}(X, Y)$, then $g \circ f \sim g' \circ f$.

(2) If $f, f' \in \text{Hom}_{\mathcal{C}}(X, Y)$ are equivalent and $g \in \text{Hom}_{\mathcal{C}}(Y, Z)$, then $g \circ f \sim g \circ f'$.

We write $[f]$ for the equivalence class of f .

Given such an equivalence relation, we can form a new category to be denoted by \mathcal{C}/\sim as follows:

□

$$\cdot \text{Ob}(\mathcal{C}/\sim) := \text{Ob}(\mathcal{C})$$

$\cdot \text{Hom}_{\mathcal{C}/\sim}(X, Y) := \text{Hom}_{\mathcal{C}}(X, Y)/\sim$ - the set of equivalence classes

$$\cdot [g] \circ [f] = [g \circ f] \text{ - well-defined precisely b/c of (1) \& (2)}$$

We note that there is a natural functor $\pi: \mathcal{C} \rightarrow \mathcal{C}/\sim$ given by $X \mapsto X, f \mapsto [f]$.

Example: Let M be a monoid. Note that the equivalence class of $1 \in M$ is a submonoid, say M_0 , moreover, (1) & (2) imply that $mM_0 = M_0m \forall m \in M$. Such submonoids are called normal (for groups we recover the usual condition). And if $M_0 = [1]$ is normal, then (1) and (2) hold - an exercise. For a normal submonoid M_0 we can M/M_0 with a natural monoid structure - just as we do for groups. The category \mathcal{C}/\sim corresponds to the quotient monoid M/M_0 and the functor π is just the natural epimorphism $M \rightarrow M/M_0$.

Rem:* \mathcal{C}/\sim looks like a quotient category. But in situations where the term "quotient" is used and that are closer to quotients of abelian groups (Serre quotients of abelian categories) the construction is different - and more difficult.

B2) Homotopy category of topological spaces.

Let's recall the usual category of topological spaces. Let

X be a set. One can define the notion of topology on X : we declare some subsets of X to be "open", these are supposed to satisfy certain axioms. A set w. topology is called a topological space. A map $f: X \rightarrow Y$ of topological spaces is called continuous if $U \subset Y$ is open $\Rightarrow f^{-1}(U) \subset X$ is open. We define the category Top of topological spaces w.

$\text{Ob}(\text{Top}) = \text{topological spaces.}$

$\text{Hom}_{\text{Top}}(X, Y) := \text{continuous maps } X \rightarrow Y$

Composition = composition of maps.

One issue: this category is hard to understand - hard to study topological spaces up to homeomorphisms.

Now we introduce our equivalence relation of $\text{Hom}_{\text{Top}}(X, Y)$

Definition: Continuous maps $f_0, f_1: X \rightarrow Y$ are called **homotopic** if \exists a continuous map $F: X \times [0, 1] \rightarrow Y$ s.t. $f_0(x) = F(x, 0)$ & $f_1(x) = F(x, 1)$.

Informally, f_0, f_1 are homotopic if one can continuously deform f_0 to f_1 . It turns out that being homotopic is an equivalence relation satisfying (1) & (2) from B1. The corresponding category Top/\sim is known as the **homotopy category of topol'l spaces**. Note that in this category morphisms are not maps!

9 | Here is why we care about the homotopy category. Isomorphic

here means homotopic (X is homotopic to Y if $\exists X \xrightarrow{f} Y$, $Y \xrightarrow{g} X$ s.t. fg is homotopic to 1_Y & gf is homotopic to 1_X) and this is easier to understand than being homeomorphic. Second, the classical invariants such as homology and homotopy groups only depend on homotopy type. A more educated way to state this: these invariants are functors from the homotopy category of topological spaces to Groups (true as stated for homology, for homotopy it's more subtle, this requires fixing a point in X and hence need to work w. an auxiliary category of "pointed" topological spaces - up to homotopy).