

Lecture 19: Categories, functors & functor morphisms IV.

1) Coproducts.

2) Adjoint functors.

BONUS: Adjunction unit & counit.

Refs: [R], Section 4.1; [HS], Sections II.5, II.7

1) Coproducts.

Let $\mathcal{C}$ be a category.

Definition: Let $X_1, X_2 \in \text{Ob}(\mathcal{C})$. Their **coproduct** (that we denote by $X_1 * X_2$) is the product in $\mathcal{C}^{\text{opp}}$ i.e.

(I) $F_{X_1 * X_2} \xrightarrow{\cong} F_{X_1} \times F_{X_2}$, where we write F_X for the Hom functor $\text{Hom}_{\mathcal{C}}(X, \cdot): \mathcal{C} \rightarrow \text{Sets}$.

(II) equivalently, there are morphisms $X_i \xrightarrow{\iota_i} X_1 * X_2$, $i=1,2$, s.t. $\forall Y \in \text{Ob}(\mathcal{C})$ & $X_i \xrightarrow{\varphi_i} Y$, $i=1,2$, $\exists! \varphi: X_1 * X_2 \rightarrow Y \mid \varphi_i = \varphi \circ \iota_i$

The equivalence of (I) & (II) follows from Lemma in Sec 2 of Lec 18 (where we replace $\mathcal{C}$ w. $\mathcal{C}^{\text{opp}}$).

Examples: 1) Let $\mathcal{C} = \text{Sets}$. Then $X_1 * X_2 = X_1 \sqcup X_2$ (and ι_i is the natural inclusion). (II) is manifest.

2) Let $\mathcal{C} = A\text{-mod}$. Then $X_1 * X_2 = X_1 \oplus X_2$: for any A -module Y , have a natural isomorphism

$$\eta_Y: \text{Hom}_A(X_1 \oplus X_2, Y) \xrightarrow{\sim} \text{Hom}_A(X_1, Y) \times \text{Hom}_A(X_2, Y)$$

see Sec 1 of Lec 4. To check (η_Y) is a functor morphism is an exercise.

Later on we will describe the coproduct in the category of commutative A -algebras (this will be the tensor product).

2) Adjoint functors.

Let $\mathcal{C}, \mathcal{D}$ be categories. Being "adjoint" is the most important relationship that a functor $\mathcal{C} \rightarrow \mathcal{D}$ can have with a functor $\mathcal{D} \rightarrow \mathcal{C}$.

2.1) Definition

Let $F: \mathcal{C} \rightarrow \mathcal{D}, G: \mathcal{D} \rightarrow \mathcal{C}$ be functors.

Definition: F is left adjoint to G (and G is right adjoint to F) if:

$$\forall X \in \text{Ob}(\mathcal{C}), Y \in \text{Ob}(\mathcal{D}) \exists \text{ bijection } \eta_{X,Y}: \text{Hom}_{\mathcal{D}}(F(X), Y) \xrightarrow{\sim} \text{Hom}_{\mathcal{C}}(X, G(Y)) \text{ s.t.}$$

$$(1) \forall X, X' \in \text{Ob}(\mathcal{C}), Y \in \text{Ob}(\mathcal{D}), X' \xrightarrow{\varphi} X \quad (\leadsto F(X') \xrightarrow{F(\varphi)} F(X))$$

the following is commutative:

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$$\begin{array}{ccc}
 \text{Hom}_{\mathcal{D}}(F(X), Y) & \xrightarrow{\cong} & \text{Hom}_{\mathcal{C}}(X, G(Y)) \\
 \downarrow ? \circ F(\varphi) & & \downarrow ? \circ G \\
 \text{Hom}_{\mathcal{D}}(F(X'), Y) & \xrightarrow{\cong} & \text{Hom}_{\mathcal{C}}(X', G(Y))
 \end{array}$$

(2) $\forall Y, Y' \in \text{Ob}(\mathcal{D}), Y \xrightarrow{\psi} Y', X \in \text{Ob}(\mathcal{C})$, the following is comm'ive

$$\begin{array}{ccc}
 \text{Hom}_{\mathcal{D}}(F(X), Y) & \xrightarrow{\cong} & \text{Hom}_{\mathcal{C}}(X, G(Y)) \\
 \downarrow \psi \circ ? & & \downarrow G(\psi) \circ ? \\
 \text{Hom}_{\mathcal{D}}(F(X), Y') & \xrightarrow{\cong} & \text{Hom}_{\mathcal{C}}(X, G(Y'))
 \end{array}$$

For us the main reason to consider adjoint functors is that we can get interesting functors as adjoints to boring (e.g. forgetful) functors.

2.2) Examples.

Below A is a commutative ring.

Example 1: Let G be $\text{For}: A\text{-Mod} \rightarrow \text{Sets}$ (forgetful)
 $F := \text{Free}: \text{Sets} \rightarrow A\text{-Mod}$ (Example 4a in Sec 1.2 of Lec 17),
 $\text{Free}(I) = A^{\oplus I}$ & for $f: I \rightarrow J$ (map of sets): $\text{Free}(f)(e_i) = e_{f(i)}$.

Claim: F is left adjoint to G

Below we write Maps for Hom_{Sets} (& Hom_A for $\text{Hom}_{A\text{-Mod}}$)

• construct $\eta_{I,M}: \text{Hom}_A(A^{\oplus I}, M) \xrightarrow{\sim} \text{Maps}(I, M)$

$$\begin{array}{ccc} \text{Hom}_A(A^{\oplus I}, M) & \xrightarrow{\sim} & \text{Maps}(I, M) \\ \downarrow \tau & \searrow \eta_{I,M} & \downarrow \psi \\ & & [i \mapsto \tau(e_i)] \end{array}$$

basis element

• check commutative diagram (1): $\forall$ maps $\varphi: I \rightarrow J$

$$\begin{array}{ccc} \tau \in \text{Hom}_A(A^{\oplus J}, M) & \xrightarrow[\sim]{\eta_{J,M}} & \text{Maps}(J, M) \\ \downarrow ? \circ \text{Free}(\varphi) & & \downarrow ? \circ \varphi \\ \text{Hom}_A(A^{\oplus I}, M) & \xrightarrow[\sim]{\eta_{I,M}} & \text{Maps}(I, M) \\ \downarrow & & \downarrow \\ \tau \mapsto [\text{unique } \tau': A^{\oplus I} \rightarrow M \text{ s.t. } \tau'(e_i) := \tau(e_{\varphi(i)})] & & \\ \downarrow & & \downarrow \\ \tau \mapsto [i \mapsto \tau(e_{\varphi(i)})] & \xleftarrow{\quad} & \\ \downarrow & & \downarrow \\ \tau \mapsto [j \mapsto \tau(e_j)] & \xleftarrow{\quad} & \end{array}$$

Check (2): $\forall \psi \in \text{Hom}_A(M, N)$, the following is commutative

$$\begin{array}{ccc} \tau \in \text{Hom}_A(A^{\oplus I}, M) & \xrightarrow[\sim]{\eta_{I,M}} & \text{Maps}(I, M) \\ \downarrow \psi \circ ? & & \downarrow \psi \circ ? \text{ where now } \psi \text{ is viewed as map of sets} \\ \text{Hom}_A(A^{\oplus I}, N) & \xrightarrow[\sim]{\eta_{I,N}} & \text{Maps}(I, N) \end{array}$$

Both $\downarrow$ & $\downarrow$ send τ to $[i \mapsto \psi(\tau(i))]$.

The adjunction is established

Example 2: Let $S \subset A$ be a multiplicative subset $\leadsto$ localization $A[S^{-1}]$ w. ring homomorphism $\iota: A \rightarrow A[S^{-1}]$. So we get functors $F := \cdot[S^{-1}]: A\text{-Mod} \rightarrow A[S^{-1}]\text{-Mod}$ and $G := \iota^*: A[S^{-1}]\text{-Mod} \rightarrow A\text{-Mod}$ (pullback = forgetful functor).

Claim: F is left adjoint to G .

For $M \in \text{Ob}(A\text{-Mod})$, $N \in \text{Ob}(A[S^{-1}]\text{-Mod})$, we have a bijection

$$\eta_{M,N}: \text{Hom}_{A[S^{-1}]}(M[S^{-1}], N) \xrightarrow{\sim} \text{Hom}_A(M, N) \quad (\text{we omit } \iota^* \text{ from the notation})$$

$$\begin{array}{ccc} \downarrow \psi & \xrightarrow{\quad} & \downarrow \psi \circ \iota_M \\ \psi & \xrightarrow{\quad} & \psi \circ \iota_M \end{array}$$
 (where $\iota_M: M \rightarrow M[S^{-1}]$, $m \mapsto \frac{m}{1}$), this is 4) of Proposition in Sec 2.2 of Lec 9.

Now we need to show that diagrams (1) and (2) from Sec 2.1 commute. Let's check (1): for $\tau \in \text{Hom}_A(M_1, M_2)$ need to show

$$\begin{array}{ccc} \text{Hom}_{A[S^{-1}]}(M_2[S^{-1}], N) & \xrightarrow{? \circ \iota_{M_2}} & \text{Hom}_A(M_2, N) \\ \downarrow ? \circ \tau[S^{-1}] & & \downarrow ? \circ \tau \quad \text{commutes} \\ \text{Hom}_{A[S^{-1}]}(M_1[S^{-1}], N) & \xrightarrow{? \circ \iota_{M_1}} & \text{Hom}_A(M_1, N) \end{array}$$

$\rightarrow \downarrow$ gives $? \circ \overbrace{\iota_{M_2} \circ \tau}^{M_1 \rightarrow M_2[S^{-1}]}$, and $\downarrow \rightarrow$ gives $? \circ \overbrace{\tau[S^{-1}] \circ \iota_{M_1}}^{M_1 \rightarrow M_2[S^{-1}]}$; for $m \in M_1$, have $\iota_{M_2} \circ \tau(m) = \frac{\tau(m)}{1}$, $\tau[S^{-1}] \circ \iota_{M_1}(m) = \tau[S^{-1}]\left(\frac{m}{1}\right) = \frac{\tau(m)}{1}$. So $\iota_{M_2} \circ \tau = \tau[S^{-1}] \circ \iota_{M_1}$ and the diagram indeed commutes.

Diagram (2) becomes: for $\zeta \in \text{Hom}_{A[S^{-1}]}(N_1, N_2)$:

$$\begin{array}{ccc}
 \text{Hom}_{A[S^{-1}]}(M[S^{-1}], N_1) & \xrightarrow{\eta_{M, N_1} = ? \circ \iota_M} & \text{Hom}_A(M, N_1) \\
 \downarrow \zeta \circ ? & & \downarrow \zeta \circ ? \quad [\iota^*(\zeta) = \zeta \text{ b/c } \iota^* \text{ is forgetful}] \\
 \text{Hom}_{A[S^{-1}]}(M[S^{-1}], N_1) & \xrightarrow{\eta_{M, N_2} = ? \circ \iota_M} & \text{Hom}_A(M, N_2)
 \end{array}$$

It is commutative.

2.3) Uniqueness.

Proposition: If $F', F^2: \mathcal{C} \rightarrow \mathcal{D}$ are left adjoint to $G: \mathcal{D} \rightarrow \mathcal{C}$, then $F^2 \cong F'$.

Proof: Suppose we have $\eta_{X,Y}^i: \text{Hom}_{\mathcal{D}}(F^i(X), Y) \xrightarrow{\sim} \text{Hom}_{\mathcal{C}}(X, G(Y))$ that make (1) & (2) commutative.

$\eta_{X,Y} := (\eta_{X,Y}^2)^{-1} \circ \eta_{X,Y}^1: \text{Hom}_{\mathcal{D}}(F^1(X), Y) \xrightarrow{\sim} \text{Hom}_{\mathcal{D}}(F^2(X), Y)$ that make the following analogs of (1) and (2) commutative (exercise).

(1) $\forall X' \xrightarrow{\varphi} X:$

$$\begin{array}{ccc}
 \text{Hom}_{\mathcal{D}}(F^1(X), Y) & \xrightarrow{\eta_{X,Y}^1 \quad ? \circ \tau_X} & \text{Hom}_{\mathcal{D}}(F^2(X), Y) \\
 \downarrow ? \circ F^1(\varphi) & & \downarrow ? \circ F^2(\varphi) \\
 \text{Hom}_{\mathcal{D}}(F^1(X'), Y) & \xrightarrow{\eta_{X',Y}^1 \quad ? \circ \tau_{X'}} & \text{Hom}_{\mathcal{D}}(F^2(X'), Y)
 \end{array}$$

(2) $\forall Y \xrightarrow{\psi} Y':$

$$\begin{array}{ccc}
 \text{Hom}_{\mathcal{D}}(F^1(X), Y) & \xrightarrow{\eta_{X,Y}^1 \quad ? \circ \tau_X} & \text{Hom}_{\mathcal{D}}(F^2(X), Y) \\
 \downarrow \psi \circ ? & & \downarrow \psi \circ ? \\
 \text{Hom}_{\mathcal{D}}(F^1(X), Y') & \xrightarrow{\eta_{X,Y'}^1 \quad ? \circ \tau_X} & \text{Hom}_{\mathcal{D}}(F^2(X), Y')
 \end{array}$$

Fix X , look at (2): it tells us that $\eta_{X,\cdot}$ is a functor morphism (and hence isomorphism - b/c each $\eta_{X,Y}$ is bijection) between $\text{Hom}_{\mathcal{D}}(F^1(X), \cdot)$ & $\text{Hom}_{\mathcal{D}}(F^2(X), \cdot)$. By Yoneda Lemma, have the unique isomorphism $\tau_X \in \text{Hom}_{\mathcal{D}}(F^2(X), F^1(X))$ s.t.
 $\eta_{X,Y} = ? \circ \tau_X$. Plug this into diagrams (1) & (2).

We now show that τ is a functor morphism $F^2 \Rightarrow F^1$ (hence an isom'm b/c each τ_X is an iso): we need to show the diagram

$$(*) \quad \begin{array}{ccc} F^2(X') & \xrightarrow{\tau_{X'}} & F^1(X') \\ \downarrow F^2(\varphi) & & \downarrow F^1(\varphi) \\ F^2(X) & \xrightarrow{\tau_X} & F^1(X) \end{array}$$

is commutative. Indeed, (1) is commutative, so

$$\psi \circ (\tau_X \circ F^2(\varphi)) = \psi \circ (F^1(\varphi) \circ \tau_{X'}) \quad \forall Y \in \text{Ob}(\mathcal{D}), \psi \in \text{Hom}_{\mathcal{D}}(F^1(X), Y).$$

Take $Y = F^1(X)$, $\psi = 1_{F^1(X)}$ & get that $(*)$ is commutative. $\square$

2.4) Remarks.

1) Fix X & consider composition of functors

$$\mathcal{D} \xrightarrow{G} \mathcal{C} \xrightarrow{\text{Hom}_{\mathcal{C}}(X, \cdot)} \text{Sets}$$

If F is left adj't to G , then $F(X)$ represents this composition via isomorphism $\eta_{X,\cdot}$, see Diagram (2) in Sec 2.1

2) We can view $\text{Hom}_{\mathcal{C}}(\cdot, ?)$ as a functor $\mathcal{C}^{\text{opp}} \times \mathcal{C} \rightarrow \text{Sets}$

Similarly for $\mathcal{D}$ compositions $\mathcal{C}^{op} \times \mathcal{D} \rightarrow \text{Sets}$

$$\text{Hom}_{\mathcal{D}}(F(\cdot), ?), \text{Hom}_{\mathcal{C}}(\cdot, G(?))$$

Diagrams (1) & (2) combine to show that $[F \text{ is left adj't to } G] \Leftrightarrow$ the two functors above are isomorphic (via η, ϵ)

3) Many categorical notions (including adjunction) have parallels in Linear Algebra. Let $\mathbb{F}$ be a field. There's a distinguished vector space, $\mathbb{F}$. For a finite dimensional vector space V , we can consider its dual, V^* . Have a vector space pairing $\langle \cdot, \cdot \rangle: V^* \times V \rightarrow \mathbb{F}$, $\langle \alpha, v \rangle := \alpha(v)$. And for a linear map $A: V \rightarrow W$ we can consider its **adjoint**, the unique linear map $A^*: W^* \rightarrow V^*$ s.t. $\langle \beta, Av \rangle = \langle A^* \beta, v \rangle$.

Here are analogs of this for categories. An analog of $\mathbb{F}$ is Sets . An analog of passing from V to V^* is passing from a category $\mathcal{C}$ to the category $\mathcal{C}^{op}$. An analog of linear maps $U \rightarrow V$ is functors $\mathcal{C} \rightarrow \mathcal{D}$. An analog of the pairing $V^* \times V \rightarrow \mathbb{F}$ is $\text{Hom}_{\mathcal{C}}(\cdot, \cdot): \mathcal{C}^{op} \times \mathcal{C} \rightarrow \text{Sets}$. Finally an analog of $\langle A^* \beta, v \rangle = \langle \beta, Av \rangle$ is our definition of adjoint functors.

There are differences as well. First, a functor $\mathcal{C} \rightarrow \mathcal{D}$ is the same thing as a functor $\mathcal{C}^{op} \rightarrow \mathcal{D}^{op}$ but there's no way to get a linear map $V^* \rightarrow W^*$ from $V \rightarrow W$. Also adjunction of functors is very sensitive to the sides (the left adjoint of G may not be isomorphic to the right adjoint - moreover exactly one of those may fail to exist), while for linear maps this issue doesn't arise.

BONUS: adjunction unit & counit.

Let $F: \mathcal{C} \rightarrow \mathcal{D}$ be left adjoint to $G: \mathcal{D} \rightarrow \mathcal{C}$. We claim that this gives rise to functor morphisms: the adjunction unit $\varepsilon: \text{Id}_{\mathcal{C}} \Rightarrow GF$ & counit $\eta: FG \Rightarrow \text{Id}_{\mathcal{D}}$.

We construct ε and leave η as an exercise.

Consider $X_1, X_2 \in \text{Ob}(\mathcal{C})$. Then we have the bijection

$$\eta_{X_1, F(X_2)}: \text{Hom}_{\mathcal{D}}(F(X_1), F(X_2)) \xrightarrow{\sim} \text{Hom}_{\mathcal{C}}(X_1, GF(X_2))$$

Note that F gives rise to a map $\text{Hom}_{\mathcal{C}}(X_1, X_2) \rightarrow \text{Hom}_{\mathcal{D}}(F(X_1), F(X_2))$

Composing this map w. the bijection $\eta_{X_1, F(X_2)}$ we get

$$\varepsilon_{X_1, X_2}: \text{Hom}_{\mathcal{C}}(X_1, X_2) \rightarrow \text{Hom}_{\mathcal{C}}(X_1, GF(X_2)).$$

Now we can argue as in the proof of Proposition 1.3 to see that $\exists! \varepsilon: \text{Id}_{\mathcal{C}} \Rightarrow GF$ s.t. $\varepsilon_{X_1, X_2}(\varphi) = \varepsilon_{X_2} \circ \varphi$.

A natural question to ask is: for two functors $F: \mathcal{C} \rightarrow \mathcal{D}$, $G: \mathcal{D} \rightarrow \mathcal{C}$ & functor morphisms $\varepsilon: \text{Id}_{\mathcal{C}} \Rightarrow GF$, $\eta: FG \Rightarrow \text{Id}_{\mathcal{D}}$ when is F left adjoint to G (& ε, η unit & counit).

Very Premium Exercise: TFAE

- F is left adjoint to G w. unit ε & counit η
- The composed morphisms $F \Rightarrow FGF \Rightarrow F$, $G \Rightarrow GFG \Rightarrow G$ induced by ε, η (cf. Problem 8 in HW3) are the identity endomorphisms (of F & G).