Ruantizations of affine schemes

(addendum to Lec 22).

1) Correspondence between truncated quantizations. Let X be a finite type affine Poisson scheme, and A=C[X]. For K71, we can talk about Kth truncated quantizations of A and of Q_X , Sec. 1.1 of Lec 22. It turns out, they are in bijection. Here's one (easier) direction.

Proposition 1: Let Die be a Kth truncated quantization of Ox. Then $\Gamma(\mathcal{D}_{t,\kappa})$ is a kth truncated quantin of A. Proof: A Poisson isomorphism $D_{\mu}/(h) \longrightarrow O_{\chi}$ gives rise to a Poisson embedding (Dt,)/(t) ~ A. We need to show it's an isomorphism. For this we use the long exact sequence in cohomology for the functor RT applied to the SES 0 -> DIA T DIA -> OX -> O. We use H'(X, O)=0 to deduce $H^{1}(\mathcal{D}_{h,\kappa}) = 0$ for all K. And this implies $\Gamma(\mathcal{D}_{h,\kappa}) \longrightarrow A$.

П

To get from restricted quantizations of A to those of X we need to localize. First, let's review how localization works for <u>non commutative</u> vines. Suppose B is a noncommutative ving and SCB is a subset containing 1 and closed under taking products. We want the localization B[S'] to satisfy the universal property and consist of right fractions as", REB, SES. It's known that B[S"] exists provided the following Ore conditions are satisfied:

01: Every right fraction is also a left fraction: HaeB, SES = 6EB, tES s.t. tR = 65 (<> "25"=t"6") Also, every left fraction is also a right fraction. 02: If seS, aEB are s.t. sa=0, then ItES s.t. at=0. And the other way around (for as=0).

Now we get back to our situation: B= St, a Kth truncated quantization of A. Pick feA, take its lift f $to St, and set S = {f, n, n, o}.$

Lemma: S setisfies (01) & (02) Sketch of proof: Note that [f,] = 0. To establish (01) for s= f, a EA we take t= f and a suitable 6 (commute f through a K times). Details as well as checking (02) are left as exercise.

Exercise 1: Prove that every other lift \tilde{f} of f to $\mathcal{H}_{t,\kappa}$ gives an invertible element in $\mathcal{H}_{t,\kappa}$ [S"]. Deduce that $\mathcal{H}_{f,\kappa}[S^{-1}]$ is independent of the choice of f.

So we get an algebra to be denoted by Sty [f"]. Exercise 2: 1) Prove it's flat over C[t]/(tr). 2) Produce a natural algebra homomorphism $\mathscr{G}_{f_k}[f^{-'}] \longrightarrow \mathscr{G}_{f_k}[(f_q)^{-'}].$

Finally, we have the following result generalizing its usual _commutative analog

Proposition 2: 1) There's a unique sheat of algebras Loc (St.) on X s.t. $\cdot \quad \Gamma(X_{f}, Loc(\mathcal{G}_{f, f})) \xrightarrow{\sim} \mathcal{G}_{f, f}[f']$ • $\Gamma(X_{f}, Loc(\mathcal{A}_{f,\mu})) \longrightarrow \Gamma(X_{f}, Loc(\mathcal{A}_{f,\mu}))$ coincides w. the homomorphism of [f"] -> of [(fg)"]. 2) The sheet Loc (At) is a kth truncated quantization of Ox. 3) The maps $\mathcal{D}_{\mu,\kappa} \mapsto \Gamma(\mathcal{D}_{\mu,\kappa})$ are $\mathcal{G}_{\mu,\kappa} \mapsto Loc(\mathcal{G}_{\mu,\kappa})$ are mutually inverse bijections between the (isomorphism classes of) with truncated quantizations of X and of A.

2) Correspondence between formal quantizations. To a formal quantization Dy on Ox we assign its global sections ((D)). This is a formal quantization of A= C(X). Indeed, arguing as in the proof of Prop. 1 in Sec 2, we see that $\Gamma(\mathcal{D}_{h,\kappa_1})/(h^{\kappa}) \xrightarrow{\sim} \Gamma(\mathcal{D}_{h,\kappa})$. Then one uses $\Gamma(\mathcal{D}_{h,\kappa})$ = (lim Dt, k) = lim (Dt, k) to conclude that (Dt) 15 c formal quantization of A. Details are exercise.

Now suppose we are given a formal quantization St, of A. Let A = A ((h") Pick f A. The universal property of localization yields a homomorphism $\mathcal{R}_{h_{K+1}}[f^{-\prime}] \rightarrow \mathcal{R}_{h_{K}}[f^{-\prime}].$

Exercise: This homomorphism induces an isomorphism $\mathscr{H}_{\mathsf{f}_{\mathsf{K}+1}}[f^{-\prime}]/(f^{\mathsf{k}}) \xrightarrow{\sim} \mathscr{H}_{\mathsf{f}_{\mathsf{K}}}[f^{-\prime}].$ Moreover, these homomorphisms glue together to $Loc(\mathcal{G}_{t,\kappa})/(\hbar^{\kappa}) \xrightarrow{\sim} Loc(\mathcal{G}_{t,\kappa})$ (*)

Then we can define Loc(Sty) as lim Loc(Sty). It's a formal quantization of \mathcal{O}_{χ} ; $(*) \Rightarrow Loc(\mathfrak{S}_{\chi})/(\mathfrak{h}^{\kappa}) \xrightarrow{\sim} Loc(\mathfrak{S}_{\chi})$ $\forall \kappa$, these claims are left as exercises.

On the other hand, have $D_{j} = \lim_{k \to \infty} D_{jk}$, hence $\Gamma(D_{j}) =$ $\lim_{k \to \infty} \Gamma(\mathcal{D}_{h^{k}}). \text{ We have } \Gamma(\mathcal{D}_{h^{K+1}})/(h^{k}) \xrightarrow{\sim} \Gamma(\mathcal{D}_{h^{k}}), \forall K, \text{ hence}$ $\Gamma(\mathcal{D}_{tx}) = \Gamma(\mathcal{D}_{t})/(h^{*})$. Now Proposition 2 from Sec 1 implies that the assignments $\mathcal{D}_{t} \mapsto \Gamma(\mathcal{D}_{t}) \& \mathscr{A}_{t} \mapsto Loc(\mathscr{A}_{t})$ are mutually inverse to each other.